

A Convex Relaxation Framework for Strategic Bidding in Electricity Markets

Mahdi Ghamkhari
Department of Computer Science
University of California Davis

Outline

M. Ghamkhari, A. Sadeghi-Mobarakeh, H. Mohsenian-Rad, “Strategic Bidding for Producers in Nodal Electricity Markets: A Convex Relaxation Approach,”
Accepted for Publication in IEEE Transactions on Power Systems, July 2016

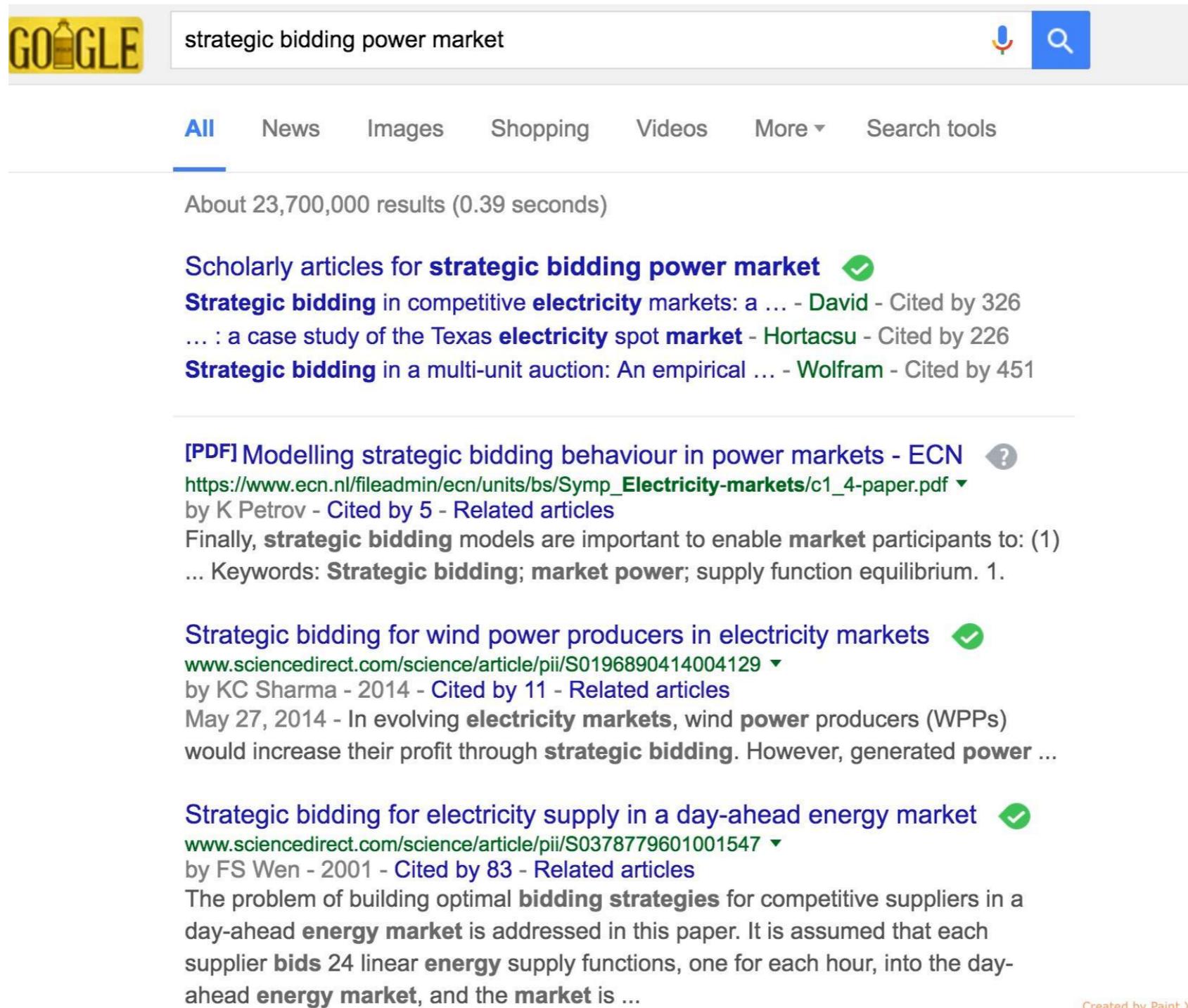
Joint work with Ashkan Sadeghi-Mobarakeh and Hamed Mohsenian-Rad

History

[Strategic gaming analysis for electric power systems: an MPEC ...](#) 
ieeexplore.ieee.org/iel5/59/18773/00867153.pdf ▾
by BF Hobbs - 2000 - [Cited by 661](#) - [Related articles](#)
... **MPEC** Approach. Benjamin F. **Hobbs**, Member, IEEE, Carolyn B. Metzler, and Jong-Shi Pang limits the number of **bidding strategies** that can be considered.

Created by Paint X

History



A screenshot of a Google search results page. The search query "strategic bidding power market" is entered in the search bar. The results are filtered under the "All" tab, showing approximately 23,700,000 results found in 0.39 seconds. The first result is a PDF titled "Modelling strategic bidding behaviour in power markets - ECN" from the European Energy Research Alliance (ECN). It discusses strategic bidding models and their importance for market participants. The second result is an article from ScienceDirect titled "Strategic bidding for wind power producers in electricity markets" by KC Sharma, published in 2014. It explores how wind power producers can increase their profit through strategic bidding. The third result is an article from ScienceDirect titled "Strategic bidding for electricity supply in a day-ahead energy market" by FS Wen, published in 2001. It addresses the problem of building optimal bidding strategies for competitive suppliers in a day-ahead energy market.

strategic bidding power market

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About 23,700,000 results (0.39 seconds)

Scholarly articles for **strategic bidding power market** ✓

Strategic bidding in competitive **electricity** markets: a ... - David - Cited by 326

... : a case study of the Texas **electricity** spot **market** - Hortacsu - Cited by 226

Strategic bidding in a multi-unit auction: An empirical ... - Wolfram - Cited by 451

[PDF] Modelling strategic bidding behaviour in power markets - ECN ?

https://www.ecn.nl/fileadmin/ecn/units/bs/Symp_Electricity-markets/c1_4-paper.pdf ▾

by K Petrov - Cited by 5 - Related articles

Finally, **strategic bidding** models are important to enable **market** participants to: (1) ... Keywords: **Strategic bidding**; **market power**; supply function equilibrium. 1.

Strategic bidding for wind power producers in electricity markets ✓

www.sciencedirect.com/science/article/pii/S0196890414004129 ▾

by KC Sharma - 2014 - Cited by 11 - Related articles

May 27, 2014 - In evolving **electricity** markets, wind power producers (WPPs) would increase their profit through **strategic bidding**. However, generated power ...

Strategic bidding for electricity supply in a day-ahead energy market ✓

www.sciencedirect.com/science/article/pii/S0378779601001547 ▾

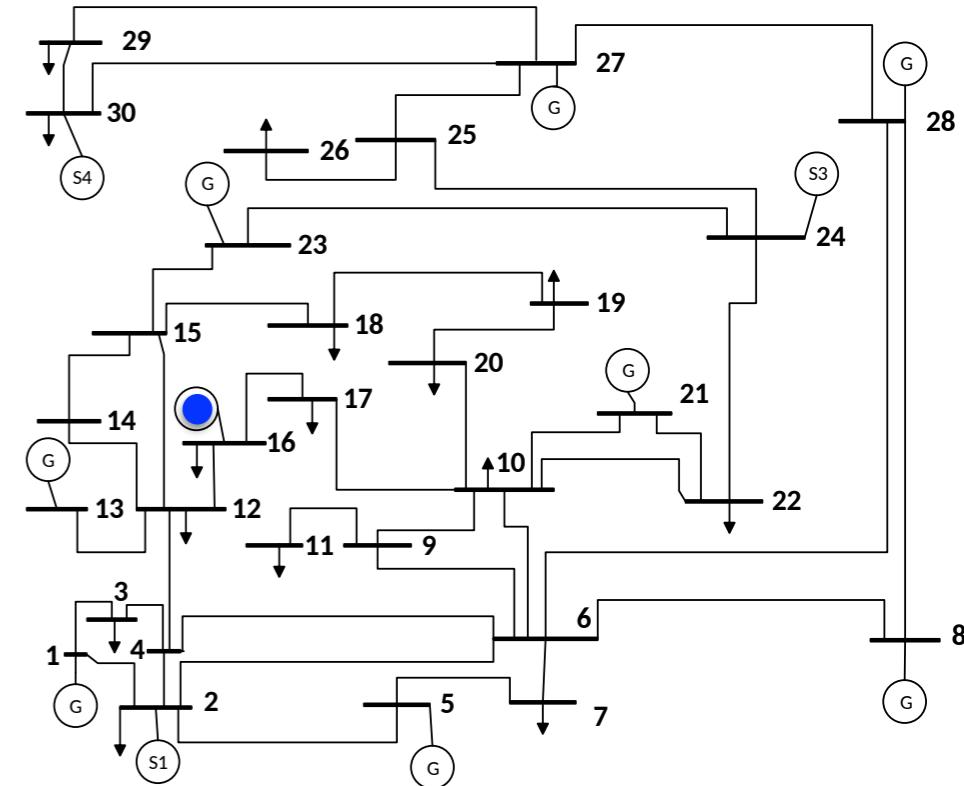
by FS Wen - 2001 - Cited by 83 - Related articles

The problem of building optimal **bidding strategies** for competitive suppliers in a day-ahead **energy market** is addressed in this paper. It is assumed that each supplier **bids** 24 linear **energy** supply functions, one for each hour, into the day-ahead **energy market**, and the **market** is ...

Electricity Market



Consumers



Electricity Network constitutes of
Generators, Consumers and
Transmission Lines

- Strategic Generator seeks to maximizes its profit by bidding in a strategic way

MPEC

**Mathematical Program with Equilibrium Constraints
(MPEC)**

$$\begin{aligned} & \text{Maximize } x^T F x + 2 f^T x \\ & p_i^T x + p_{i0} \geq 0 \quad \forall i \\ & v_m^T x + v_{m0} = 0 \quad \forall m \\ & x^T Q_z x + 2 q_z^T x = 0 \quad \forall z \end{aligned}$$

$$Q_z = d_z q_z^T$$

Will be
needed later

Inherent relation
between parameters

Mixed Integer Linear Program

$$0 \leq q_z^T x \leq \text{Binary} \times (\text{LargeNumber})$$
$$0 \leq d_z^T x + 2 \leq 1 - \text{Binary} \times (\text{LargeNumber})$$

Binary Variable

Binary Variable

$$\begin{aligned} & \text{Maximize } x^T F x + 2 f^T x \\ & p_i^T x + p_{i0} \geq 0 \quad \forall i \\ & v_m^T x + v_{m0} = 0 \quad \forall m \\ \Leftrightarrow & x^T Q_z^T x + 2 q_z^T x = 0 \quad \forall z \end{aligned}$$

$$Q_z = d_z q_z^T$$

Pool Strategy of a Producer With Endogenous Formation of Locational ... 

ieeexplore.ieee.org/iel5/59/5282387/05272241.pdf

by C Ruiz - 2009 - Cited by 136 - Related articles

Index Terms—Electricity pool, endogenous price formation., LMP, offering strategy ... bought, and the hourly locational marginal prices (LMPs). Specifically, this ...

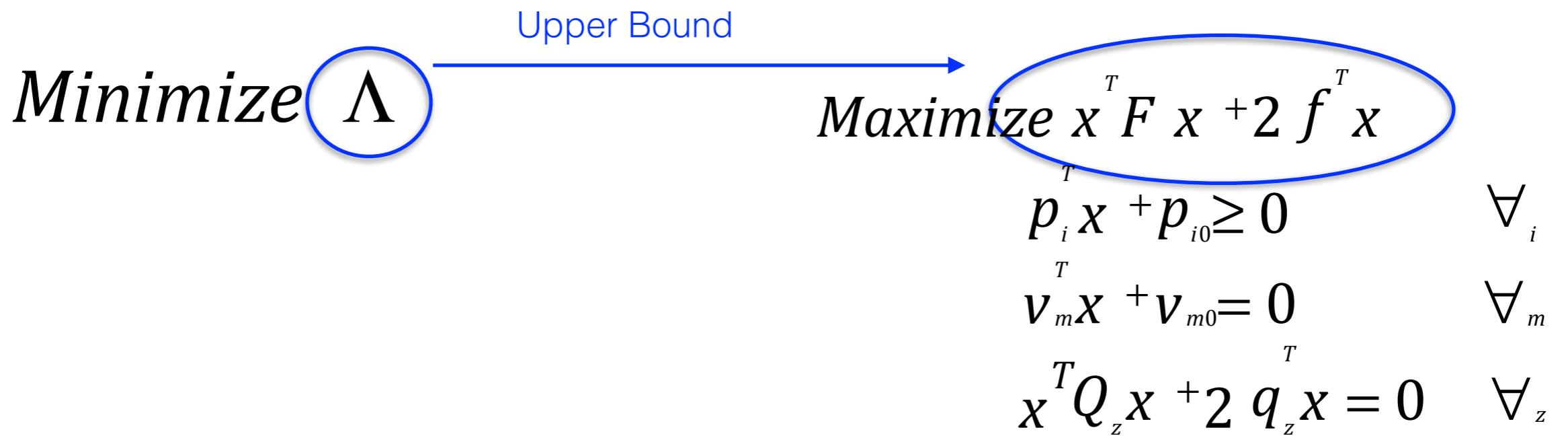
Solutions

- MILP: gives global solution
- MILP: Computation time increases Exponentially

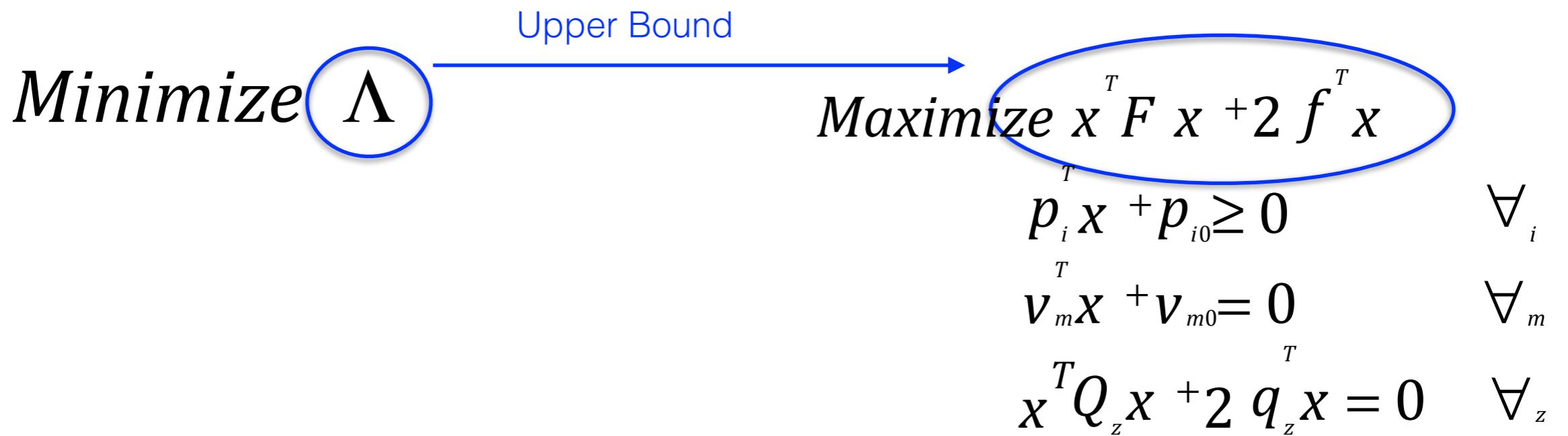
Solutions

- MILP: gives global solution
- Our Approach: gives global solution with 99% Optimality
- MILP: Computation time increases Exponentially
- Our Approach:Computation Time increases Linearly

Our Approach



Our Approach



Λ is upper bound if and only if

$$\Lambda - x^T F x - 2 f^T x \geq 0$$

is positive on

$$\left\{ \begin{array}{ll} p_i^T x + p_{i0} \geq 0 & \forall i \\ v_m^T x + v_{m0} = 0 & \forall m \\ x^T Q_z x + 2 q_z^T x = 0 & \forall z \end{array} \right\}$$

Positivstellensatz

Polynomials that are positive on semi Algebraic sets

$$\Lambda - x^T F x - 2 f^T x \geq 0$$

is positive on

$$\left\{ \begin{array}{l} p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ v_m^T x + v_{m0} = 0 \quad \forall_m \\ x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \end{array} \right.$$

Polynomial

Semi Algebraic Set

Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

⋮

$$\sum_{m=1}^M \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

$$\Lambda - x^T F x - 2 f^T x \geq 0 \quad \text{is positive on}$$

$$\left\{ \begin{array}{l} p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ v_m^T x + v_{m0} = 0 \quad \forall_m \\ x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \end{array} \right\}$$

Polynomial

Semi Algebraic Set

Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I (\text{SOS Polynomial}) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I (\text{SOS Polynomial}) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I (\text{SOS Polynomial}) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

:

$$\sum_{m=1}^M (\text{Arbitrary Polynomial}) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z (\text{Arbitrary Polynomial}) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

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Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

:

$$\sum_{m=1}^M \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

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Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

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⋮

$$\sum_{m=1}^M \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

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Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

⋮

$$\sum_{m=1}^M \left(\begin{array}{c} Arbitrary \\ Polynomial \end{array} \right) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} Arbitrary \\ Polynomial \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

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Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I \left(\begin{array}{c} \text{SOS} \\ \text{Polynomial} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

⋮

$$\sum_{m=1}^M \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (V_m^T x + V_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} \text{Arbitrary} \\ \text{Polynomial} \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

$$\Lambda - x^T F x - 2 f^T x \geq 0 \quad \text{is positive on}$$

$$\left\{ \begin{array}{l} p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ V_m^T x + V_{m0} = 0 \quad \forall_m \\ x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \end{array} \right.$$

Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

⋮

$$\sum_{m=1}^M \left(\begin{array}{c} Arbitrary \\ Polynomial \end{array} \right) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} Arbitrary \\ Polynomial \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

$$\Lambda - x^T F x - 2 f^T x \geq 0 \quad \text{is positive on}$$

$$\left\{ \begin{array}{l} p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ v_m^T x + v_{m0} = 0 \quad \forall_m \\ x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \end{array} \right.$$

Schmudgen Positivstellensatz

$$\begin{aligned}
 & \Lambda - x^T F x - 2 f^T x - \\
 & \sum_{i=1}^I (\text{SOS Polynomial}) (p_i^T x + p_{i0}) - \\
 & \sum_{i=1}^I \sum_{j=1}^I (\text{SOS Polynomial}) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) - \\
 & \sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I (\text{SOS Polynomial}) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) - \\
 & \vdots \\
 & \sum_{m=1}^M (\text{Arbitrary Polynomial}) (v_m^T x + v_{m0}) - \\
 & \sum_{z=1}^Z (\text{Arbitrary Polynomial}) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n
 \end{aligned}$$

$$\Lambda - x^T F x - 2 f^T x \geq 0$$

is positive on

$$\left\{
 \begin{array}{ll}
 p_i^T x + p_{i0} \geq 0 & \forall_i \\
 v_m^T x + v_{m0} = 0 & \forall_m \\
 x^T Q_z x + 2 q_z^T x = 0 & \forall_z
 \end{array}
 \right.$$

Variables are polynomials

Schmudgen Positivstellensatz

$$\begin{aligned}
 & \Lambda - x^T F x - 2 f^T x - \\
 & \sum_{i=1}^I (\text{SOS Polynomial})(p_i^T x + p_{i0}) - \\
 & \sum_{i=1}^I \sum_{j=1}^I (\text{SOS Polynomial})(p_i^T x + p_{i0})(p_j^T x + p_{j0}) - \\
 & \sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I (\text{SOS Polynomial})(p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) - \\
 & \vdots \\
 & \sum_{m=1}^M (\text{Arbitrary Polynomial})(v_m^T x + v_{m0}) - \\
 & \sum_{z=1}^Z (\text{Arbitrary Polynomial})(x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n
 \end{aligned}$$

Minimize Λ

$$\Lambda - x^T F x - 2 f^T x \geq 0 \quad \text{is positive on}$$

$$\left\{
 \begin{array}{ll}
 p_i^T x + p_{i0} \geq 0 & \forall_i \\
 v_m^T x + v_{m0} = 0 & \forall_m \\
 x^T Q_z x + 2 q_z^T x = 0 & \forall_z
 \end{array}
 \right.$$

Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \underset{\text{Polynomial}}{\underset{\text{SOS}}{(p_i^T x + p_{i0}) -}}$$

$$\sum_{i=1}^I \sum_{j=1}^I \underset{\text{Polynomial}}{\underset{\text{SOS}}{(p_i^T x + p_{i0})(p_j^T x + p_{j0}) -}}$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I \underset{\text{Polynomial}}{\underset{\text{SOS}}{(p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -}}$$

⋮

$$\sum_{m=1}^M \underset{\text{Arbitrary Polynomial}}{\underset{\text{SOS}}{(v_m^T x + v_{m0}) -}} \xrightarrow{\text{Arbitrary Polynomials}}$$

$$\sum_{z=1}^Z \underset{\text{Arbitrary Polynomial}}{\underset{\text{SOS}}{(x^T Q_z^T x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n}}$$

$$\Lambda - x^T F x - 2 f^T x \geq 0 \quad \text{is positive on}$$

$$\left\{ \begin{array}{l} p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ v_m^T x + v_{m0} = 0 \quad \forall_m \\ x^T Q_z^T x + 2 q_z^T x = 0 \quad \forall_z \end{array} \right\}$$

Schmudgen Positivstellensatz

$$\Lambda - x^T F x - 2 f^T x -$$

$$\sum_{i=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I \left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0})(p_t^T x + p_{t0}) -$$

:

$$\sum_{m=1}^M \left(\begin{array}{c} Arbitrary \\ Polynomial \end{array} \right) (v_m^T x + v_{m0}) -$$

$$\sum_{z=1}^Z \left(\begin{array}{c} Arbitrary \\ Polynomial \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n$$

$$\left(\begin{array}{c} SOS \\ Polynomial \end{array} \right) = \sum (Polynomial)^2$$

x is not variables
x is an index for infinite constraints

$$\Lambda - x^T F x - 2 f^T x \geq 0 \quad \text{is positive on}$$

$$\left\{ \begin{array}{l} p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ v_m^T x + v_{m0} = 0 \quad \forall_m \\ x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \end{array} \right.$$

A Convex Optimization Problem

Minimize Λ

$$\Lambda - \mathbf{x}^T F \mathbf{x} - 2 f^T \mathbf{x} -$$

$$\sum_{i=1}^I (\text{SOS}_{\text{Polynomial}})(p_i^T \mathbf{x} + p_{i0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I (\text{SOS}_{\text{Polynomial}})(p_i^T \mathbf{x} + p_{i0})(p_j^T \mathbf{x} + p_{j0}) -$$

$$\sum_{i=1}^I \sum_{j=1}^I \sum_{t=1}^I (\text{SOS}_{\text{Polynomial}})(p_i^T \mathbf{x} + p_{i0})(p_j^T \mathbf{x} + p_{j0})(p_t^T \mathbf{x} + p_{t0}) -$$

\vdots

$$\sum_{m=1}^M (\text{Arbitrary}_{\text{Polynomial}})(V_m^T \mathbf{x} + V_{m0}) -$$

$$\sum_{z=1}^Z (\text{Arbitrary}_{\text{Polynomial}})(\mathbf{x}^T Q_z \mathbf{x} + 2 q_z^T \mathbf{x}) \geq 0 \quad \forall \mathbf{x} \in \mathbb{R}^n$$

Relaxation of Polynomials in Psatz

$$\begin{aligned} & \Lambda - x^T F x - 2 f^T x - \\ & \sum_{i=1}^I \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) (p_i^T x + p_{i0}) - \\ & \sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) - \\ & \sum_{m=1}^M \left(\begin{array}{c} \text{Linear} \\ \text{Polynomial} \end{array} \right) (V_m^T x + V_{m0}) - \\ & \sum_{z=1}^Z \left(\begin{array}{c} \text{Arbitrary} \\ \text{Scalar} \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n \end{aligned}$$



Quadratic Expression

Computationally Tractable Reformulation

$$\begin{aligned}
 & \Lambda - x^T F x - 2 f^T x - \\
 & \sum_{i=1}^I \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) (p_i^T x + p_{i0}) - \\
 & \sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) (p_i^T x + p_{i0})(p_j^T x + p_{j0}) - \\
 & \sum_{m=1}^M \left(\begin{array}{c} \text{Linear} \\ \text{Polynomial} \end{array} \right) (v_m^T x + v_{m0}) - \\
 & \sum_{z=1}^Z \left(\begin{array}{c} \text{Arbitrary} \\ \text{Scalar} \end{array} \right) (x^T Q_z x + 2 q_z^T x) \geq 0 \quad \forall x \in \mathbb{R}^n
 \end{aligned}$$

Quadratic Expression

$$\begin{bmatrix} 1 \\ x \end{bmatrix}^T \gamma \begin{bmatrix} 1 \\ x \end{bmatrix} \geq 0$$

$$\begin{aligned}
 \gamma = & \begin{pmatrix} \Lambda & f^T \\ -f & F \end{pmatrix} - \\
 & \sum_{i=1}^I \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) \begin{pmatrix} p_{i0} & \frac{p_i^T}{2} \\ \frac{p_i}{2} & \mathbf{0} \end{pmatrix} - \\
 & \sum_{i=1}^I \sum_{j=1}^I \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) \begin{bmatrix} p_{j0} \\ \frac{p_j}{2} \end{bmatrix} \begin{bmatrix} p_{j0} \\ \frac{p_j}{2} \end{bmatrix}^T - \\
 & \sum_{m=1}^M \text{Scalar} \begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} v_{m0} \\ \frac{v_m}{2} \end{bmatrix}^T - \sum_{m=1}^M \sum_{l=1}^n \text{Scalar} \begin{bmatrix} 0 \\ e_l \end{bmatrix} \begin{bmatrix} v_{m0} \\ \frac{v_m}{2} \end{bmatrix}^T \\
 & \sum_{z=1}^Z \left(\begin{array}{c} \text{Positive} \\ \text{Scalar} \end{array} \right) \begin{pmatrix} 0 & \frac{q_i^T}{2} \\ \frac{q_i}{2} & Q_z \end{pmatrix}
 \end{aligned}$$

Computationally Tractable

$$\begin{bmatrix} 1 \\ x \end{bmatrix}^T \gamma \begin{bmatrix} 1 \\ x \end{bmatrix} \geq 0 \quad \iff \quad \gamma \succeq 0$$

Semi Definite

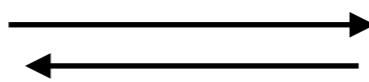
Computationally Tractable

$$\Lambda$$

is upper bound if

$$\gamma \succeq 0$$

$$\begin{bmatrix} 1 \\ x \end{bmatrix}^T \gamma \begin{bmatrix} 1 \\ x \end{bmatrix} \geq 0$$



$$\gamma \succeq 0$$

Semi Definite

Computationally Tractable

Minimize Λ

$$\gamma \succeq 0$$

Best *customized* upper bound

$$\text{Maximize } x^T F x + 2 f^T x$$

$$p_i^T x + p_{i0} \geq 0 \quad \forall_i$$

$$v_m^T x + v_{m0} = 0 \quad \forall_m$$

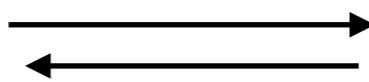
$$x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z$$

$$\Lambda$$

is upper bound if

$$\gamma \succeq 0$$

$$\begin{bmatrix} 1 \\ x \end{bmatrix}^T \gamma \begin{bmatrix} 1 \\ x \end{bmatrix} \geq 0$$



$$\gamma \succeq 0$$

Semi Definite

Computationally Tractable

Minimize Λ

$$\gamma \succeq 0$$



Λ is 97% optimal

Computationally Tractable

Minimize Λ

$$\gamma \succeq 0$$



Λ is 97% optimal



What happened to x ?

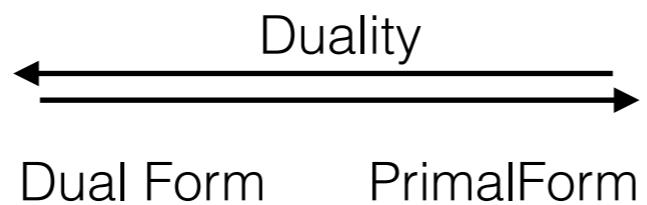
x : Optimal optimization variables in MPEC ?

$$\begin{bmatrix} 1 \\ x \end{bmatrix}^T \begin{bmatrix} 1 \\ x \end{bmatrix} \geq 0 \iff \gamma \succeq 0$$

Recovery

Minimize Λ

$$\gamma \succeq 0$$



$$\text{maximize } \text{trace} \begin{pmatrix} 0 & f^T \\ f & F \end{pmatrix} X$$

$$\text{trace} \begin{pmatrix} p_{i0} & \frac{p_i^T}{2} \\ \frac{p_i}{2} & 0 \end{pmatrix} X \geq 0 \quad \forall i$$

$$\text{trace} \begin{pmatrix} e_l \begin{bmatrix} v_{m0} \\ v_m \end{bmatrix}^T & X \end{pmatrix} = 0 \quad \forall m, l$$

$$\text{trace} \begin{pmatrix} \begin{bmatrix} p_{i0} \\ p_i \end{bmatrix} \begin{bmatrix} p_{j0} \\ p_j \end{bmatrix}^T & X \end{pmatrix} \geq 0 \quad \forall i, j$$

$$\text{trace} \begin{pmatrix} 0 & \frac{q_z^T}{2} \\ \frac{q_z}{2} & Q_z \end{pmatrix} X \geq 0 \quad \forall z$$

$$X_{11} = 1$$

$$X \succeq 0$$

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

First column of X

$$\text{maximize } \text{trace} \begin{pmatrix} 0 & f^T \\ f & F \end{pmatrix} X$$

$$\text{trace} \begin{pmatrix} p_{i0} & \frac{p_i}{2}^T \\ \frac{p_i}{2} & 0 \end{pmatrix} X \geq 0 \quad \forall i$$

$$\text{trace} \begin{pmatrix} e_l \begin{bmatrix} v_{m0} \\ v_m \end{bmatrix}^T & X \end{pmatrix} = 0 \quad \forall m, l$$

$$\text{trace} \begin{pmatrix} \begin{bmatrix} p_{i0} \\ p_i \end{bmatrix} \begin{bmatrix} p_{j0} \\ p_j \end{bmatrix}^T & X \end{pmatrix} \geq 0 \quad \forall i, j$$

$$\text{trace} \begin{pmatrix} \begin{pmatrix} 0 & q_z^T \\ q_z & Q_z \end{pmatrix} & X \end{pmatrix} \geq 0 \quad \forall z$$

$$X_{11} = 1$$

$$X \succeq 0$$

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

$$\text{maximize } \text{trace} \begin{pmatrix} 0 & f^T \\ f & F \end{pmatrix} X$$

$$\text{trace} \begin{pmatrix} p_{i0} & \frac{p_i}{2}^T \\ \frac{p_i}{2} & 0 \end{pmatrix} X \geq 0 \quad \forall i$$

$$\text{trace} \begin{pmatrix} e_l \begin{bmatrix} v_{m0} \\ v_m \end{bmatrix}^T & X \end{pmatrix} = 0 \quad \forall m, l$$

$$\text{trace} \begin{pmatrix} \begin{bmatrix} p_{i0} \\ p_i \end{bmatrix} \begin{bmatrix} p_{j0} \\ p_j \end{bmatrix}^T & X \end{pmatrix} \geq 0 \quad \forall i, j$$

$$\text{trace} \begin{pmatrix} \begin{bmatrix} 0 & q_z^T \\ q_z & Q_z \end{bmatrix} & X \end{pmatrix} \geq 0 \quad \forall z$$

$$X_{11} = 1$$

$$X \succeq 0$$

First column of X



X^* is not feasible in MPEC

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}^T$$

we use X^* to produce a feasible solution to MPEC

$$\begin{aligned} & \text{Maximize } x^T F x + 2 f^T x \\ & p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ & v_m^T x + v_{m0} = 0 \quad \forall_m \\ & x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \\ & Q_z = d_z q_z^T \end{aligned}$$

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}^T$$

* we use X^* to produce a feasible solution to MPEC

$$\begin{aligned} & \text{Maximize } x^T F x + 2 f^T x \\ & p_i^T x + p_{i0} \geq 0 \quad \forall_i \\ & v_m^T x + v_{m0} = 0 \quad \forall_m \\ & x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z \end{aligned}$$

$\xleftarrow{\hspace{10em}}$

$$Q_z = d_z q_z^T$$

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}^T$$

*
we use X^* to produce a feasible solution to MPEC

$$|d_z^T x^* + 2| \leq \varepsilon \quad \text{and} \quad |q_z^T x^*| \geq \Delta \implies (d_z^T x + 2) = 0$$

$$(d_z^T x + 2)(q_z^T x) = 0$$

A diagram showing two arrows originating from the term $d_z^T x + 2$. One arrow points to the term $q_z^T x$, and another arrow points to the term 0 .

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}^T$$

*
we use X^* to produce a feasible solution to MPEC

$$\begin{aligned} & (d_z^T x + 2) (q_z^T x) = 0 \\ & |d_z^T x + 2| \geq \Delta \quad \text{and} \quad |q_z^T x| \leq \varepsilon \implies (q_z^T x) = 0 \end{aligned}$$

Recovery

Approximated solution
for variables in MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}^T$$

* we use X^* to produce a feasible solution to MPEC

$$\text{Maximize } x^T F x + 2 f^T x$$

$$p_i^T x + p_{i0} \geq 0 \quad \forall_i$$

$$v_m^T x + v_{m0} = 0 \quad \forall_m$$

$$x^T Q_z x + 2 q_z^T x = 0 \quad \forall_z$$

Algorithm

1. Solve the dual of customized Psatz Relaxation

$$\begin{aligned}
 & \text{maximize} \text{ trace} \left(\begin{pmatrix} 0 & f^T \\ f & F \end{pmatrix} X \right) \\
 & \text{trace} \left(\begin{pmatrix} \frac{p_{i0}}{2} & \frac{p_i}{2}^T \\ \frac{p_i}{2} & 0 \end{pmatrix} X \right) \geq 0 \quad \forall i \\
 & \text{trace} \left(e_i \begin{bmatrix} v_{m0} \\ v_m \end{bmatrix}^T X \right) = 0 \quad \forall m, l \\
 & \text{trace} \left(\begin{bmatrix} p_{i0} \\ p_i \end{bmatrix} \begin{bmatrix} p_{j0} \\ p_j \end{bmatrix}^T X \right) \geq 0 \quad \forall i, j \\
 & \text{trace} \left(\begin{pmatrix} 0 & q_z^T \\ q_z & Q_z \end{pmatrix} X \right) \geq 0 \quad \forall z \\
 & X_{11} = 1 \\
 & X \succeq 0
 \end{aligned}$$

2. Obtain an approximated solution for MPEC

$$\begin{bmatrix} 1 \\ x^* \end{bmatrix} = X^* \begin{bmatrix} 1 \\ 0 \end{bmatrix}^T$$

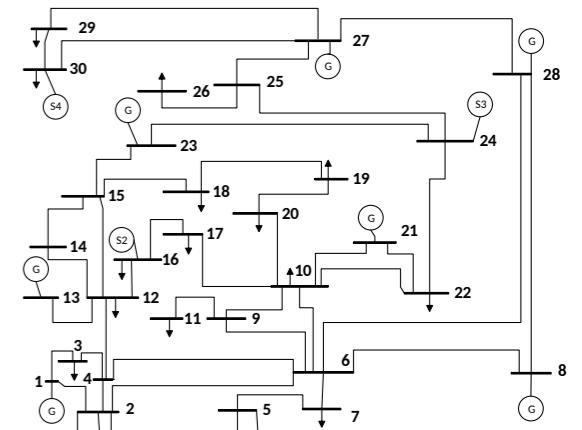
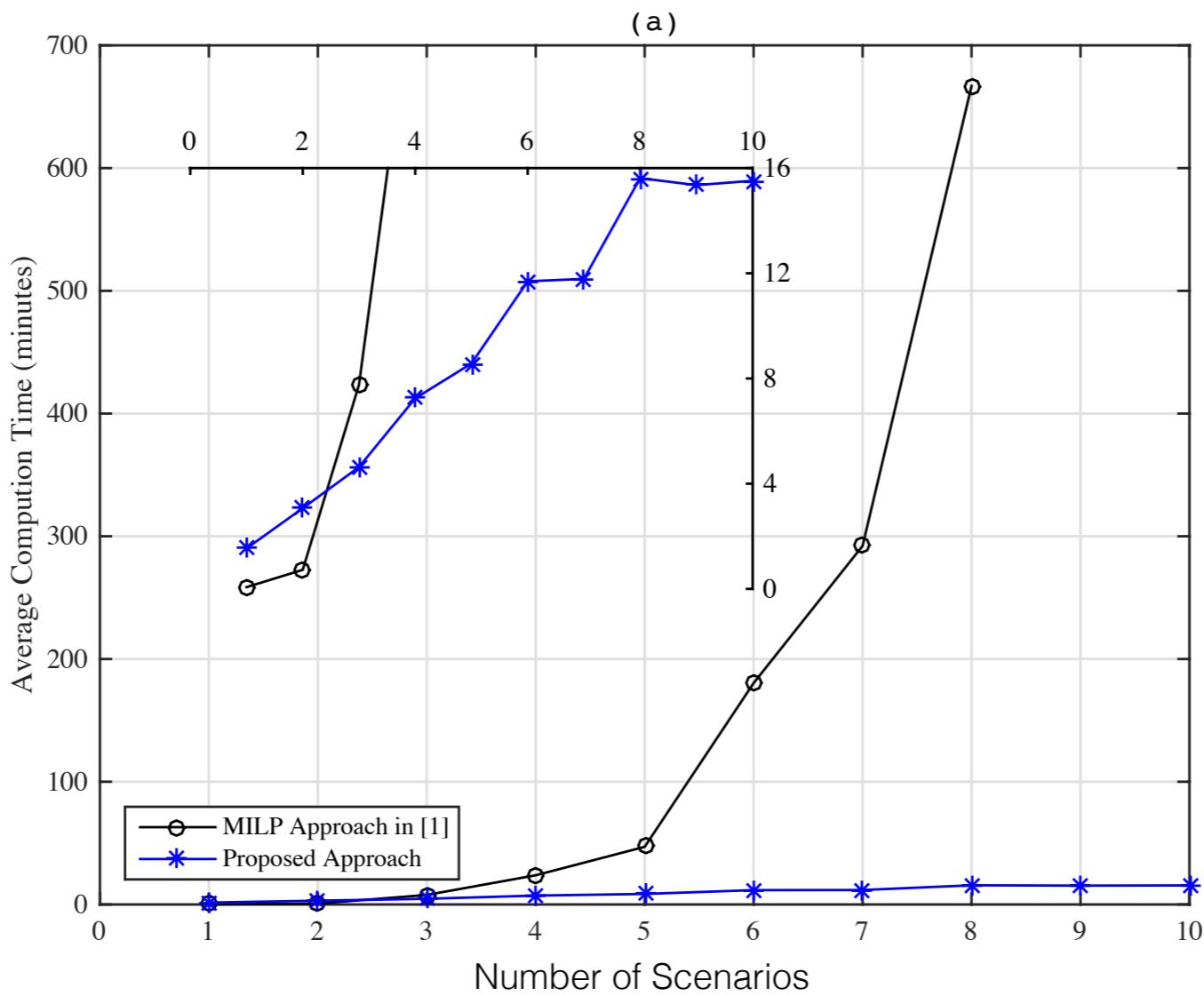
3. Eliminate some non-convex constraints

$$\begin{aligned}
 (d_z^T x + 2)(q_z^T x) &= 0 \\
 (d_z^T x + 2) &\rightarrow (d_z^T x + 2) = 0 \\
 (q_z^T x) &\rightarrow (q_z^T x) = 0
 \end{aligned}$$

4. Solve a reduced complexity MILP

$$\begin{aligned}
 & \text{Maximize } x^T F x + 2 f^T x \\
 & p_i^T x + p_{i0} \geq 0 \quad \forall i \\
 & v_m^T x + v_{m0} = 0 \quad \forall m \\
 & x^T Q_z x + 2 q_z^T x = 0 \quad \forall z
 \end{aligned}$$

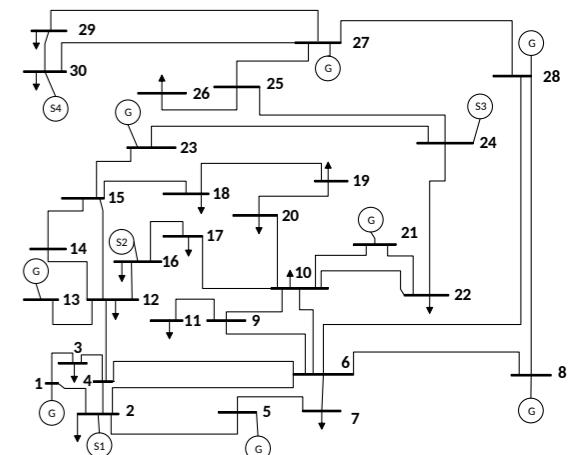
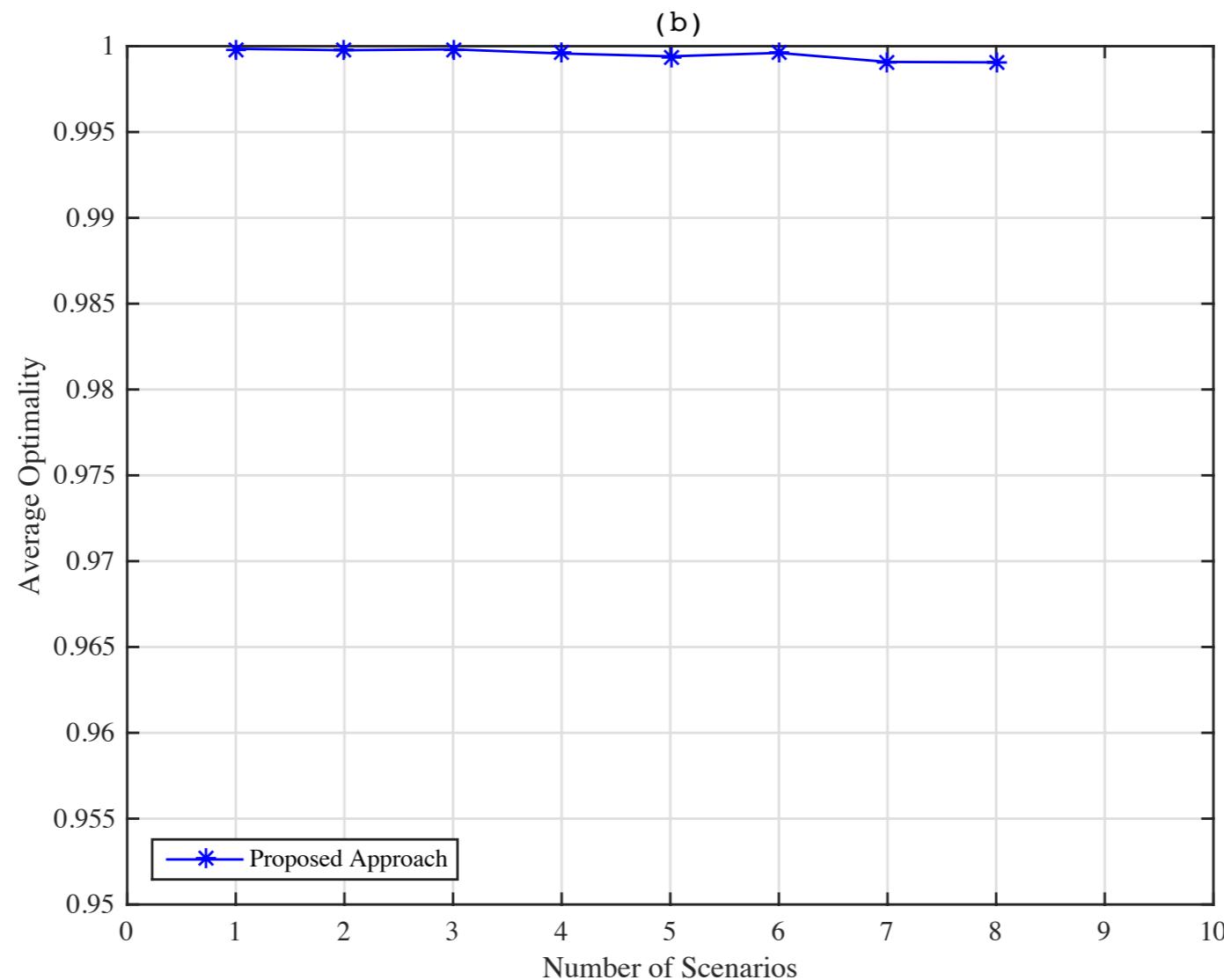
Simulation Results



IEEE 30-Bus Network

The impact of increasing the number of random scenarios on the computation time of the proposed approach and the MILP approach

Simulation Results



IEEE 30-Bus Network

The impact of increasing the number of random scenarios on the optimality of the proposed approach.

Thank You